Programming Parallel Computers

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Part 2C: How to benefit from vector operations?

With a normal scalar hammer, it does not matter much where your nails are I can do this!

With a normal scalar hammer, it does not matter much where your nails are One by one!

With a normal *scalar hammer*, it does not matter much where your nails are



With a normal *scalar hammer*, it does not matter much where your nails are



With a normal *scalar hammer*, it does not matter much where your nails are

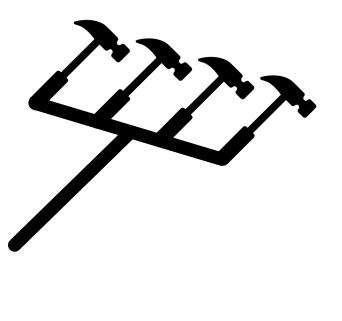


With a normal *scalar hammer*, it does not matter much where your nails are

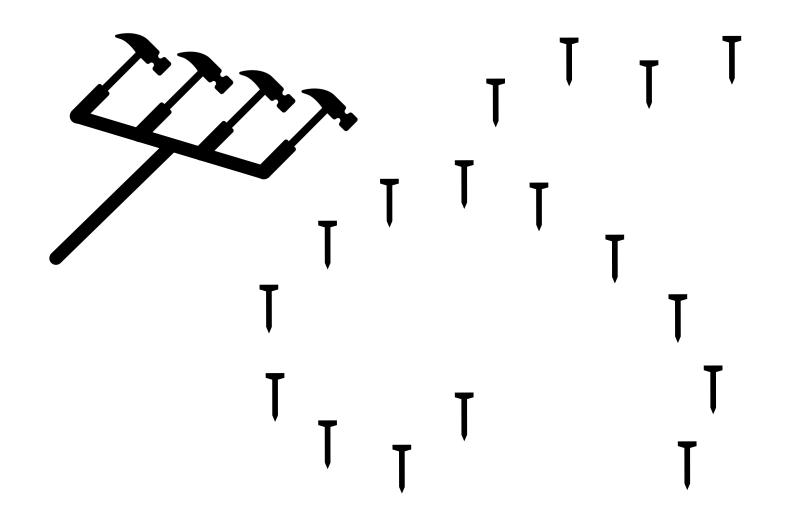
But can I do this faster?

Then you get a brand-new **vector hammer**!

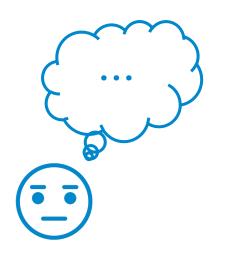


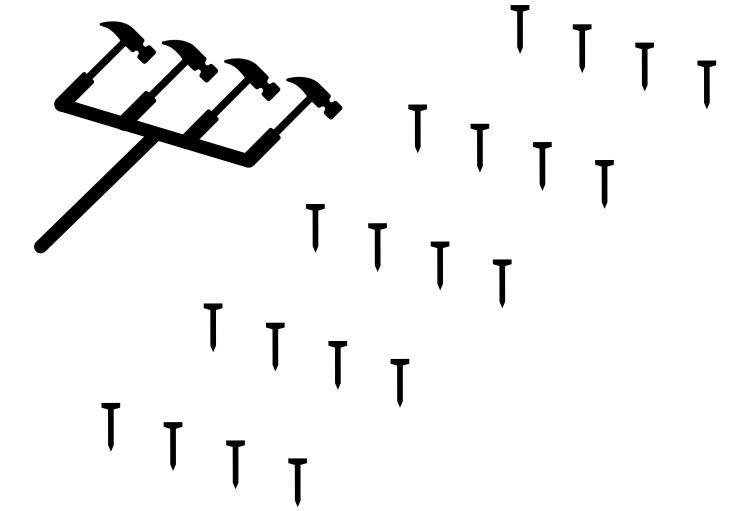


But it does not seem to make any sense to use it in your project?

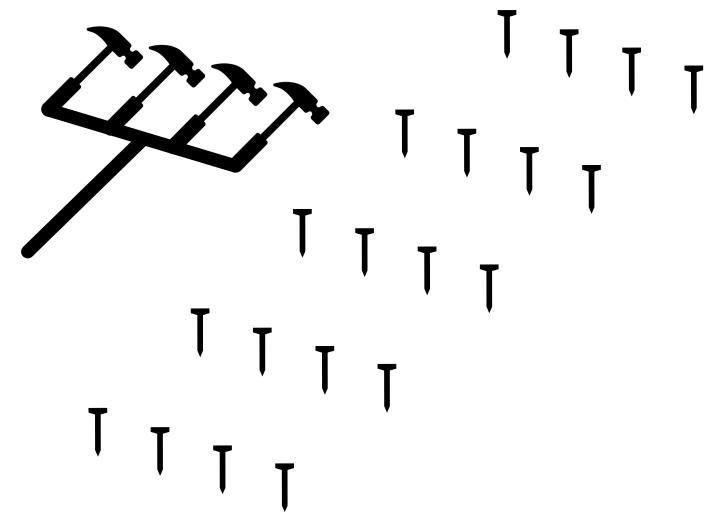


Redesign your project keeping in mind that you are wielding a vector hammer!





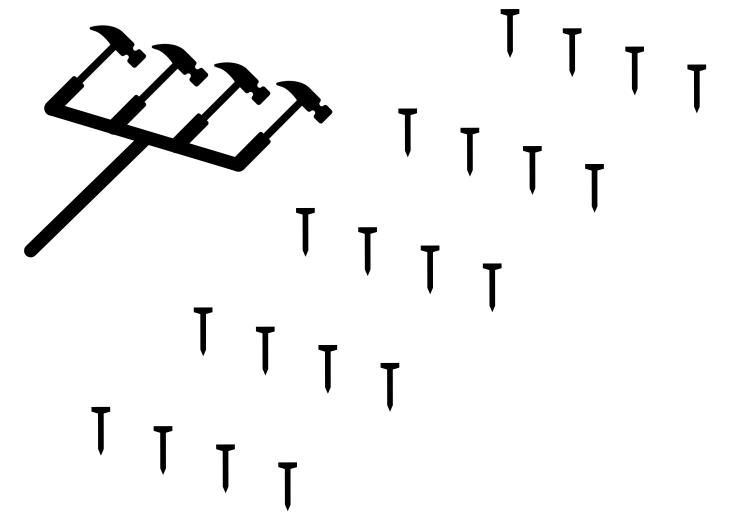
You will often have extra steps in your program to rearrange data so that inner loops can do lots of useful work with vector operations



hammer Form follows <u>function</u>

Sometimes you will need to re-think the entire data layout

Plenty of room for creativity!



- Typical idea:
 - preprocess your data
 - "pack" individual data elements to vectors
 - add *padding* if input size not multiple of 4, 8, etc.
 - do vector operations
 - "unpack" results from vectors
 - if needed, do some post-processing to turn vector results into normal results
- Make sure you do enough arithmetic operations so that all the extra work is worth it!

- Packing data, some examples:
 - vector = multiple elements from the same row of input
 - vector = one element from each row of input
 - vector = (R, G, B) triple in image processing
 - vector = one sample from each input channel in audio processing
 - vector = 256 pixels of a monochromatic image
 - vector = 32 characters of text
- Make sure you are mostly doing similar operations for each vector element
 - e.g. elementwise addition, elementwise multiplication

```
for (int i = 0; i < n; ++i) {</pre>
    for (int j = 0; j < n; ++j) {</pre>
        float v = infinity;
        for (int k = 0; k < n; ++k) {
            float x = d[n*i + k];
             float y = t[n*j + k];
            float z = x + y;
            v = min(v, z);
        r[n*i + j] = v;
```

No parallelism, scalar operations

```
for (int i = 0; i < n; ++i) {</pre>
    for (int j = 0; j < n; ++j) {</pre>
        float v0 = infinity;
float v1 = infinity;
         for (int k = 0; k < n/2; ++k) {
            float x0 = d[n*i + 2*k];
           l float x1 = d[n*i + 2*k + 1];
            float y0 = t[n*j + 2*k];
           \int float y1 = t[n*j + 2*k + 1];
            float z0 = x0 + y0;
           \int float z1 = x1 + y1;
           v0 = min(v0, z0);
v1 = min(v1, z1);
         r[n*i + j] = min(v0, v1);
```

Groups of 2 similar independent operations

```
for (int i = 0; i < n; ++i) {</pre>
    for (int j = 0; j < n; ++j) {</pre>
       float v0 = infinity;
       float v7 = infinity;
        for (int k = 0; k < n/8; ++k) {
           float x0 = d[n*i + 8*k];
           float x7 = d[n*i + 8*k + 7];
           float y0 = t[n*j + 8*k];
           float y7 = t[n*j + 8*k + 7];
           float z0 = x0 + y0;
           float z7 = x7 + y7;
           v0 = min(v0, z0);
           v7 = min(v7, z7):
        r[n*i + j] = min(v0, v1, \cdots, v7);
    }
```

Groups of 8 similar independent operations

```
for (int i = 0; i < n; ++i) {</pre>
    for (int j = 0; j < n; ++j) {</pre>
        float8_t vv = f8infty;
        for (int k = 0; k < n/8; ++k) {
            float8_t vx = vd[n/8*i + k];
            float8_t vy = vt[n/8*j + k];
            float8_t vz = vx + vy;
            vv = min8(vv, vz);
        r[n*i + j] = hmin8(vv);
```





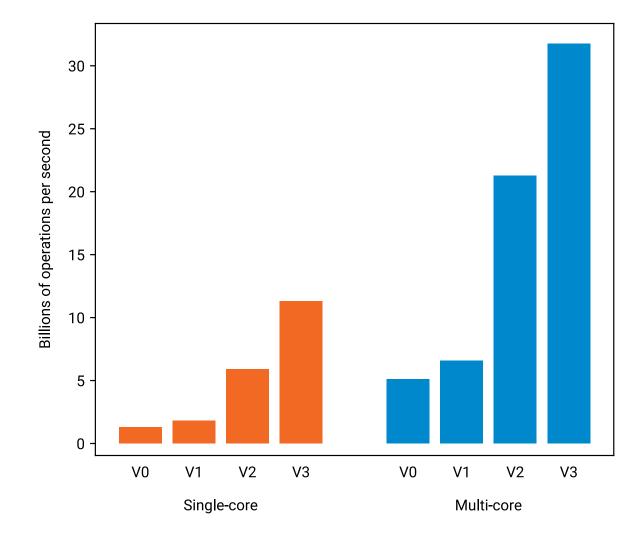
1 vector operation

Vectorization

V2: instruction-level parallelism

V3: vectorization

Running time improved from 99 s to 4 s



Data reuse will be necessary

- Performance of a typical 4-core CPU:
 - could do 64 floating-point additions per clock cycle
 - main memory bandwidth: can fetch enough data for ≈ 1.25 floating-point additions per clock cycle
 - we can only afford to fetch 2% of our input from main memory!
- Lots of data reuse needed:
 - reusing what you have got from main memory to caches
 - reusing what you have got from caches to registers
- More about this next week!