# **Programming Parallel Computers**

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Part 4C: Memory access patterns in CUDA programs

#### Sample application: cheapest 2-hop path

d (input):



d[] = { 0, 8, 2, 1, 0, 9, 4, 5, 0 }



r (output):



4, 5, 0 }

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```
void step(float* r, const float* d, int n) {
    for (int i = 0; i < n; ++i) {</pre>
        for (int j = 0; j < n; ++j) {</pre>
             float v = infinity;
             for (int k = 0; k < n; ++k) {
                 float \mathbf{x} = d[n*i + k];
                 float y = d[n*k + j];
                 float z = x + y;
                 v = min(v, z);
             r[n*i + j] = v;
```

# **Splitting work**

- Work to do:
  - need to compute *n* × *n* results
  - computing one result takes *n* steps
- How do we split this in blocks and threads?
- Natural idea:
  - one thread computes one result
  - one block computes **b** × **b** results, for some suitable b
  - if we choose e.g. b = 16, then a block consists of 8 warps

# **Splitting work**

- Example: input dimensions are 1600 x 1600:
  - we want to create 100 x 100 blocks
  - each block consists of 16 x 16 threads
- Create 10 000 blocks with 256 threads:
  - blocks numbered 0 ... 9999
  - threads numbered 0 ... 255
- Convert block & thread index to (*i*, *j*) pair:
  - thread number 123 in block number 4567 computes the result for *i* = ??? and *j* = ???

# **Splitting work**

- Example: input dimensions are 1600 x 1600:
  - we want to create 100 x 100 blocks
  - each block consists of 16 x 16 threads
- Create 10 000 blocks with 256 threads:
  - blocks numbered 0 ... 9999
  - threads numbered 0 ... 255
- Convert block & thread index to (*i*, *j*) pair:
  - thread number 123 = 7 · 16 + 11 in block number 4567 = 45 · 100 + 67 computes the result for *i* = 67 · 16 + 11 and *j* = 45 · 16 + 7

## **Splitting work: using 2D indexes**

- Example: input dimensions are 1600 x 1600:
  - we want to create 100 x 100 blocks
  - each block consists of 16 x 16 threads
- Create 10 000 blocks with 256 threads using 2D indexes:
  - blocks numbered (0, 0) ... (99, 99)
  - threads numbered (0, 0) ... (15, 15)
- Convert block & thread coordinates to (*i*, *j*) pair:
  - thread number (11, 7) in block number (67, 45)
     computes the result for *i* = 67 · 16 + 11 and *j* = 45 · 16 + 7

## **Splitting work: rounding**

- Example: input dimensions are 1601 x 1601:
  - we want to create 101 x 101 blocks
  - each block consists of 16 x 16 threads
- Create 10 000 blocks with 256 threads using 2D indexes:
  - blocks numbered (0, 0) ... (100, 100)
  - threads numbered (0, 0) ... (15, 15)
- There will be some threads with  $i \ge 1601$  and/or  $j \ge 1601$ , they will do nothing



```
float * dGPU = NULL;
cudaMalloc((void**)&dGPU, n * n * sizeof(float));
float* rGPU = NULL;
cudaMalloc((void**)&rGPU, n * n * sizeof(float));
cudaMemcpy(dGPU, d, n * n * sizeof(float),
           cudaMemcpyHostToDevice);
                                                 n/16, rounded up
dim3 dimBlock(16, 16);
dim3 dimGrid(divup(n, dimBlock.x), divup(n, dimBlock.y));
mykernel<<<dimGrid, dimBlock>>>(rGPU, dGPU, n);
cudaMemcpy(r, rGPU, n * n * sizeof(float),
           cudaMemcpyDeviceToHost);
```

```
cudaFree(dGPU); cudaFree(rGPU);
```

#### Performance

- Test input: *n* = 6300
- Maari computers:
  - baseline CPU solution: 397 s
  - best CPU solution: 2.3 s
  - current GPU solution: 42 s
- What is the bottleneck?



- A key challenge in CPU code: getting data fast enough from the CPU memory
- A key challenge in GPU code: getting data fast enough from the GPU memory

#### Memory access pattern

- Blocks are divided in warps
  - warp = 32 threads
- Entire warp executes synchronously
- If one thread reads some memory, all threads of the warp read some memory













#### Memory access pattern

- One memory read in kernel: *entire warp of threads reads memory simultaneously*
- Threads access small continuous parts of memory: need to load few cache lines  $\rightarrow \textbf{good}$
- Threads access 32 different locations far from each other: need to load many cache lines  $\rightarrow$  **bad**

First warp:

thread 0: i = 0, j = 0thread 1: **i** = 1, **j** = 0 thread 2: **i** = **2**, **j** = **0** thread 3: **i** = 3, **j** = 0 ... thread 31: **i** = **15**, **j** = **1 Pay attention** to this index!

int i = threadIdx.x + ...
int j = threadIdx.y + ...
for (... ++k) {
 float x = d[n\*i + k];
 float y = d[n\*k + j];

• • •

}

First warp, first iteration:



. . .

threads 0 & 16: read d[0] threads 1 & 17: read d[1000] threads 2 & 18: read d[2000] threads 3 & 19: read d[3000] int i = threadIdx.x + ...
int j = threadIdx.y + ...
for (... ++k) {
 float x = d[n\*i + k];
 float y = d[n\*k + j];

. . .

}

First warp, first iteration:



threads 0 & 16: read d[0] threads 1 & 17: read d[1000] threads 2 & 18: read d[2000] threads 3 & 19: read d[3000] int i = threadIdx.x + ...
int j = threadIdx.y + ...
for (... ++k) {
 float x = d[n\*i + k];
 float y = d[n\*k + j];
 ...



threads 0–15: read **d[0]** threads 16–31: read **d[1]** 

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Bad

. . .

threads 0 & 16: read d[1] threads 1 & 17: read d[1001] threads 2 & 18: read d[2001] threads 3 & 19: read d[3001]

First warp, second iteration:

int i = threadIdx.x + ...
int j = threadIdx.y + ...
for (... ++k) {
 float x = d[n\*i + k];
 float y = d[n\*k + j];
 ...



threads 0–15: read **d[1000]** threads 16–31: read **d[1001]**  First warp, third iteration:



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threads 0 & 16: read d[2] threads 1 & 17: read d[1002] threads 2 & 18: read d[2002] threads 3 & 19: read d[3002] int i = threadIdx.x + ...
int j = threadIdx.y + ...
for (... ++k) {
 float x = d[n\*i + k];
 float y = d[n\*k + j];
 ...



threads 0–15: read **d[2000]** threads 16–31: read **d[2001]** 

# Exchange the roles of i and j

int i = threadIdx.x + ...
int j = threadIdx.y + ...
for (... ++k) {
 float x = d[n\*i + k];
 float y = d[n\*k + j];
 float x = d[n\*j + k];
 float y = d[n\*k + i];

r[n\*i + j] = v;
r[n\*j + i] = v;

. . .

First warp, first iteration: int i = threadIdx.x + ... int j = threadIdx.y + ... for (... ++k) { float x = d[n\*j + k];float y = d[n\*k + i];

. . .

}



threads 0–15: read d[0] threads 16–31: read d[1000]







#### Performance

- V0: baseline 42 s
- V1: better memory access pattern 8 s
- But we can do much better by applying familiar ideas:
  - reuse data in registers
  - reuse data in "cache" (here: shared memory)

#### Performance

- V0: baseline
- V1: memory access
- V2: registers
- V3: shared memory

